Complex Event Recognition

Alexander Artikis

Elias Alevizos, Nikos Katzouris, Vagelis Michelioudakis, George Paliouras

Institute of Informatics & Telecommunications, NCSR Demokritos, Athens, Greece

http://cer.iit.demokritos.gr

Complex Event Recognition (Event Pattern Matching)

Input:

- Symbolic representation of time-stamped, simple, derived events (SDE) coming from (geographically distributed) sources.
- Big Data.

Complex Event Recognition (Event Pattern Matching)

Input:

- Symbolic representation of time-stamped, simple, derived events (SDE) coming from (geographically distributed) sources.
- Big Data.

Output:

- Complex or composite events (CE) collections of SDE and/or CE that satisfy some pattern (temporal, spatial, logical constraints).
 - Not restricted to aggregates.
- Humans understand CE easier than SDE.

Complex Event Recognition





Input	Output
340 inactive(id ₀)	
$340 \ p(id_0) = (20.88, -11.90)$	
340 appear(id ₀)	
340 walking (id_2)	
$340 p(id_2) = (25.88, -19.80)$	
340 $active(id_1)$	
$340 \ p(id_1) = (20.88, -11.90)$	
340 walking (id_3)	
$340 \ p(id_3) = (24.78, -18.77)$	
380 walking (id_3)	
$380 \ p(id_3) = (27.88, -9.90)$	
380 walking (id_2)	
$380 p(id_2) = (28.27, -9.66)$	

Input		Output
340 inactive(id ₀)	340	$left_object(id_1, id_0)$
$340 p(id_0) = (20.88, -11.90)$		
340 $appear(id_0)$		
340 walking (id_2)		
$340 p(id_2) = (25.88, -19.80)$		
340 $active(id_1)$		
$340 p(id_1) = (20.88, -11.90)$		
340 walking (id_3)		
$340 \ p(id_3) = (24.78, -18.77)$		
380 walking (id_3)		
$380 \ p(id_3) = (27.88, -9.90)$		
380 walking(id ₂)		
$380 p(id_2) = (28.27, -9.66)$		

Input		Output
340 inactive(id ₀)	340	$left_object(id_1, id_0)$
$340 \ p(id_0) = (20.88, -11.90)$	<i>since</i> (340)	$moving(id_2, id_3)$
340 $appear(id_0)$		
340 walking(id ₂)		
$340 \ p(id_2) = (25.88, -19.80)$		
340 active(id ₁)		
$340 \ p(id_1) = (20.88, -11.90)$		
340 walking(id ₃)		
$340 \ p(id_3) = (24.78, -18.77)$		
380 walking(id ₃)		
$380 \ p(id_3) = (27.88, -9.90)$		
380 walking(id ₂)		
$380 p(id_2) = (28.27, -9.66)$		





	Input	Output
200	scheduled stop enter	
215	late stop leave	
[215, 400]	abrupt acceleration	
[350, 600]	sharp turn	
620	<u>flow=low</u>	
	density=high	

	Input		Output
200	scheduled stop enter		
215	late stop leave	<i>since</i> (215)	non-punctual
[215, 400]	abrupt acceleration		
[350, 600]	sharp turn	[215, 600]	uncomfortable driving
620	<u>flow=low</u>		
	density=high	since(620)	congestion

	Input		Output
200	scheduled stop enter		
215	late stop leave	<i>since</i> (215)	non-punctual
[215, 400]	abrupt acceleration		
[350, 600]	sharp turn	[215, 600]	uncomfortable driving
620	<u>flow=low</u>		
	density=high	since(620)	congestion
700	scheduled stop enter		
720	<u>flow=low</u>		
	density=average		
820	scheduled stop leave		
915	passenger density		
	change to low		

	Input		Output
200	scheduled stop enter		
215	late stop leave	<i>since</i> (215)	non-punctual
[215, 400]	abrupt acceleration		
[350, 600]	sharp turn	[215, 600]	uncomfortable driving
620	<u>flow=low</u>		
	density=high	since(620)	congestion
700	scheduled stop enter		
720	<u>flow=low</u>		
	density=average	[620,720]	congestion
820	scheduled stop leave	[215,820]	non-punctual
915	passenger density		
	change to low		



SDE:

Credit card transactions from all over the world.

CE:

- Cloned card a credit card is being used simultaneously in different countries.
- New high use the card is being frequently used in merchants or countries never used before.
- Potential batch fraud many transactions from multiple cards in the same point-of-sale terminal in high amounts.



Fraud must be detected within 25 milliseconds.



- ► Fraud must be detected within 25 milliseconds.
- Fraudulent transactions: 0.1% of the total number of transactions.



- ► Fraud must be detected within 25 milliseconds.
- Fraudulent transactions: 0.1% of the total number of transactions.
- Fraud is constantly evolving.



- ► Fraud must be detected within 25 milliseconds.
- Fraudulent transactions: 0.1% of the total number of transactions.
- Fraud is constantly evolving.
- Erroneous transactions, missing fields.





Fast Approach



- A vessel is moving at a high speed ...
- towards other vessels.

Suspicious Delay



- ► A vessel fails to report position ...
- and the estimated speed during the communication gap is low.

Possible Rendezvous



- Two vessels are suspiciously delayed ...
- in the same location ...
- at the same time.

Package Picking



- A vessel stops in a location ...
- and another vessel stops in the same location ...
- in a short period of time.

'Sea' of information:

 400,000 vessels globally: 40,000 signals/sec to be processed online.

'Sea' of information:

- 400,000 vessels globally: 40,000 signals/sec to be processed online.
- Position signals need to be combined with other data streams
 - Weather forecasts, sea currents, etc.

'Sea' of information:

- 400,000 vessels globally: 40,000 signals/sec to be processed online.
- Position signals need to be combined with other data streams
 - Weather forecasts, sea currents, etc.
- ... and static information
 - ► NATURA areas, shallow waters, coastlines, etc.

'Sea' of noisy information:

► GPS manipulation has risen 59% over the past two years.

'Sea' of noisy information:

- ► GPS manipulation has risen 59% over the past two years.
- There is a 30% increase over the past two years of vessels reporting a false identity.

'Sea' of noisy information:

- ► GPS manipulation has risen 59% over the past two years.
- There is a 30% increase over the past two years of vessels reporting a false identity.
- ► 27% of vessels do not report position at least 10% of the time.
 - ▶ 19% of vessels are repeat offenders.

- Efficient reasoning
 - to support real-time decision-making in large-scale, (geographically) distributed applications.

- Efficient reasoning
 - to support real-time decision-making in large-scale, (geographically) distributed applications.
- Reasoning under uncertainty
 - to deal with various types of noise.

- Efficient reasoning
 - to support real-time decision-making in large-scale, (geographically) distributed applications.
- Reasoning under uncertainty
 - to deal with various types of noise.
- Automated knowledge construction
 - to avoid the time-consuming, error-prone manual CE definition development.

- Efficient reasoning
 - to support real-time decision-making in large-scale, (geographically) distributed applications.
- Reasoning under uncertainty
 - to deal with various types of noise.
- Automated knowledge construction
 - to avoid the time-consuming, error-prone manual CE definition development.

Composite Event Algebra

Core components of an event algebra with point-based semantics:

 Sequencing (SEQ) lists the required event types in temporal order — eg SEQ(A, B, C).

Composite Event Algebra

Core components of an event algebra with point-based semantics:

- Sequencing (SEQ) lists the required event types in temporal order — eg SEQ(A, B, C).
- Kleene closure (+) collects a finite yet unbound number of events of a particular type. It is used as a component of SEQ — eg SEQ(A, B+, C).
Core components of an event algebra with point-based semantics:

- Sequencing (SEQ) lists the required event types in temporal order — eg SEQ(A, B, C).
- Kleene closure (+) collects a finite yet unbound number of events of a particular type. It is used as a component of SEQ — eg SEQ(A, B+, C).
- Negation (~ or !) verifies the absence of certain events in a sequence — eg SEQ(A, !B, C).

Core components of an event algebra with point-based semantics:

- Sequencing (SEQ) lists the required event types in temporal order — eg SEQ(A, B, C).
- ▶ Kleene closure (+) collects a finite yet unbound number of events of a particular type. It is used as a component of SEQ
 — eg SEQ(A, B+, C).
- ► Negation (~ or !) verifies the absence of certain events in a sequence eg SEQ(A, !B, C).
- Value predicates specify constraints on the event attributes
 - ▶ Aggregate functions max, min, count, sum, avg.

- Composition refers to:
 - Union of constraints eg SEQ $(A, B, C) \cup$ SEQ(A, D, E).
 - Negation of a sequence eg !SEQ(A, B, C).
 - Kleene closure of a constraint eg SEQ(A, B, C)+.

- Composition refers to:
 - Union of constraints eg SEQ $(A, B, C) \cup$ SEQ(A, D, E).
 - Negation of a sequence eg !SEQ(A, B, C).
 - Kleene closure of a constraint eg SEQ(A, B, C)+.
- Windowing (WITHIN) restricts a CE definition to a specific time period.

Strict contiguity: No intervening events allowed between two sequence events in the pattern.

- Strict contiguity: No intervening events allowed between two sequence events in the pattern.
- Partition contiguity: Same as above, but the stream is partitioned into substreams according to a partition attribute. Events must be contiguous within the same partition.

- Strict contiguity: No intervening events allowed between two sequence events in the pattern.
- Partition contiguity: Same as above, but the stream is partitioned into substreams according to a partition attribute. Events must be contiguous within the same partition.
- Skip-till-next-match: Intervening events are allowed, but only non-overlapping occurrences of SEQ are detected. E.g. for SEQ(A, B, C) and a₁, b₁, b₂, c₁, only a₁, b₁, c₁ will be detected.

- Strict contiguity: No intervening events allowed between two sequence events in the pattern.
- Partition contiguity: Same as above, but the stream is partitioned into substreams according to a partition attribute. Events must be contiguous within the same partition.
- Skip-till-next-match: Intervening events are allowed, but only non-overlapping occurrences of SEQ are detected. E.g. for SEQ(A, B, C) and a₁, b₁, b₂, c₁, only a₁, b₁, c₁ will be detected.
- ► Skip-till-any-match: Most flexible (and expensive). Detects every possible occurrence. For the previous example, a₁, b₂, c₁ will also be detected.

Example

Quickly moving away from an area of suspicious activity:

- After a communication gap, ...
- a vessel changes speed to over 30 knots.
- Partition contiguity ensures that a, b, c are contiguous and refer to the same vessel (vesselld).

PATTERN SEQ(gapStart a, gapEnd b, speedChange c) WHERE partition-contiguity AND vesselld AND c.velocity > 30 WITHIN 3600

Example

Fishing pattern:

- A vessel slows down, …
- begins a series of turns, where, for each pair of successive turns, their difference in heading is more than 90 degrees, ...
- and subsequently the vessel stops moving at a low speed.

```
PATTERN SEQ(lowSpeedStart a, turn + b, lowSpeedEnd c)
WHERE skip-till-next-match
AND vesselld
AND b[i].heading -b[i-1].heading > 90
WITHIN 21600
```

CE as Chronicle

A CE can be defined as a set of events interlinked by time constraints and whose occurrence may depend on the context.

CE as Chronicle

A CE can be defined as a set of events interlinked by time constraints and whose occurrence may depend on the context.

• This is the definition of a chronicle.

CE as Chronicle

A CE can be defined as a set of events interlinked by time constraints and whose occurrence may depend on the context.

• This is the definition of a chronicle.

Chronicle recognition systems have been used in many applications:

- Cardiac monitoring system.
- Intrusion detection in computer networks.
- Distributed diagnosis of web services.

Chronicle Representation Algebra

Predicate	Meaning
event(E, T)	Event E takes place at time T
event(F:(?V1,?V2),T)	An event takes place at time T changing the value of property F from ?V1 to ?V2
<pre>noevent(E, (T1,T2))</pre>	Event E does not take place between [T1,T2)
noevent(F:(?V1,?V2), (T1,T2))	No event takes place between $[T1,T2)$ that changes the value of property F from ?V1 to ?V2
hold(F:?V, (T1,T2))	The value of property F is ?V between [T1,T2)
<pre>occurs(N,M,E,(T1,T2))</pre>	Event E takes place at least N times and at most M times between [T1,T2)

```
chronicle abnormal_vessel_movement[?id](T2) {
  event( speedChange[?id], T0 )
  event( speedChange[?id], T1 )
  event( speedChange[?id], T2 )
  T1 > T0
  T2 > T1
  T2 - T0 in [1, 20000]
  noevent( turn[?id], ( T0+1, T2 ) )
}
```

- Mathematical operators in the atemporal constraints of the language are not allowed.
 - ► No spatial reasoning.
 - No use of background knowledge.

- Mathematical operators in the atemporal constraints of the language are not allowed.
 - No spatial reasoning.
 - No use of background knowledge.
- Universal quantification is not allowed.
 - cannot express a property about all vessels in some area.

 Mathematical operators in the atemporal constraints of the language are not allowed.

- No spatial reasoning.
- No use of background knowledge.
- Universal quantification is not allowed.
 - cannot express a property about all vessels in some area.

CRS is a purely temporal reasoning system.

It is also a very efficient and scalable system.

Chronicle Recognition System: Semantics

Each CE definition is represented as a Temporal Constraint Network. Eg:



Chronicle Recognition System: Consistency Checking

Compilation stage:

- Constraint propagation in the Temporal Constraint Network.
- Consistency checking.



Chronicle Recognition System: Recognition

Recognition stage:

- ▶ Partial CE instance evolution.
- Forward (predictive) recognition.



[0,3] [1,3] В А

A: speedChange B: turn

C: stop

time



A: speedChange B: turn C: stop





A: speedChange B: turn C: stop













A: speedChange B: turn C: stop



















A@3





A: speedChange B: turn C: stop



Chronicle Recognition System

Recognition stage — partial CE instance management:

In order to manage all the partial CE instances, CRS stores them in trees, one for each CE definition.

Chronicle Recognition System

Recognition stage — partial CE instance management:

- In order to manage all the partial CE instances, CRS stores them in trees, one for each CE definition.
- Each event occurrence and each clock tick traverses these trees in order to kill some CE instances (tree nodes) or to develop some CE instances.

Chronicle Recognition System

Recognition stage — partial CE instance management:

- In order to manage all the partial CE instances, CRS stores them in trees, one for each CE definition.
- Each event occurrence and each clock tick traverses these trees in order to kill some CE instances (tree nodes) or to develop some CE instances.
- To deal with out-of-order SDE streams, CRS keeps in memory partial CE instances longer.

Chronicle Recognition System: Optimisation

Several techniques have been developed for improving efficiency. Eg, temporal focusing:

• Distinguish between rare events and frequent events.

Chronicle Recognition System: Optimisation

Several techniques have been developed for improving efficiency. Eg, temporal focusing:

- Distinguish between rare events and frequent events.
- If, according to a CE definition, a rare event should take place after a frequent event, store the incoming frequent events, and start recognition only upon the arrival of the rare events.

Chronicle Recognition System: Optimisation

Several techniques have been developed for improving efficiency. Eg, temporal focusing:

- Distinguish between rare events and frequent events.
- If, according to a CE definition, a rare event should take place after a frequent event, store the incoming frequent events, and start recognition only upon the arrival of the rare events.
- ► The number of partial CE instances is significantly reduced.
- Example:

$$(A) \xrightarrow{[1,3]} (B) \xrightarrow{[0,3]} (C) \qquad \begin{array}{c} A: \text{ speedChange} \\ B: \text{ turn} \\ C: \text{ stop} \end{array}$$
Chronicle Recognition System: Summary

- One of the earliest and most successful formal event processing systems.
- Many of its features appear in modern event processing systems.
- Very efficient and scalable event recognition.

Chronicle Recognition System: Summary

- One of the earliest and most successful formal event processing systems.
- Many of its features appear in modern event processing systems.
- Very efficient and scalable event recognition.
- ► But:
 - It is a purely temporal reasoning system.
 - It does not handle uncertainty.

Event Calculus

- A logic programming language for representing and reasoning about events and their effects.
- Key components:
 - event (typically instantaneous).
 - fluent: a property that may have different values at different points in time.

Event Calculus

- A logic programming language for representing and reasoning about events and their effects.
- Key components:
 - event (typically instantaneous).
 - fluent: a property that may have different values at different points in time.
- Built-in representation of inertia:
 - ► F = V holds at a particular time-point if F = V has been initiated by an event at some earlier time-point, and not terminated by another event in the meantime.

Run-Time Event Calculus (RTEC)

Predicate	Meaning
happensAt (E, T)	Event E occurs at time T
initiatedAt(F=V,T)	At time T a period of time for which $F = V$ is initiated
terminatedAt(F = V, T)	At time T a period of time for which $F = V$ is terminated
holdsFor(F = V, I)	I is the list of the maximal intervals for which $F = V$ holds continuously
holdsAt(F = V, T)	The value of fluent F is V at time T
union_all($[J_1, \ldots, J_n], I$)	$I = (J_1 \cup \ldots \cup J_n)$
intersect_all($[J_1, \ldots, J_n], I$)	$I = (J_1 \cap \ldots \cap J_n)$
relative_complement_all $(l', [J_1, \dots, J_n], l)$	$I=I'\setminus (J_1\cup\ldots\cup J_n)$

Run-Time Event Calculus (RTEC)

Predicate	Meaning
happensAt (E, T)	Event E occurs at time T
initiatedAt(F=V,T)	At time T a period of time for which $F = V$ is initiated
terminatedAt $(F = V, T)$	At time T a period of time for which $F = V$ is terminated
holdsFor(F = V, I)	I is the list of the maximal intervals for which $F = V$ holds continuously
holdsAt(F = V, T)	The value of fluent F is V at time T
union_all($[J_1, \ldots, J_n], I$)	$I = (J_1 \cup \ldots \cup J_n)$
intersect_all($[J_1, \ldots, J_n], I$)	$I = (J_1 \cap \ldots \cap J_n)$
relative_complement_all $(I', [J_1, \ldots, J_n], I)$	$I=I'\setminus (J_1\cup\ldots\cup J_n)$

CE definition:

. . .

 $\begin{array}{ll} \mbox{initiatedAt}(\textit{CE}, \ \textit{T}) \leftarrow \\ \mbox{happensAt}(\textit{E}_{\textit{ln}_1}, \ \textit{T}), \\ [\mbox{conditions}] \end{array}$

initiatedAt(CE, T) \leftarrow happensAt(E_{In_i}, T), [conditions] $\begin{array}{ll} \textbf{terminatedAt}(\textit{CE}, \ \textit{T}) \leftarrow \\ \textbf{happensAt}(\textit{E}_{\textit{T}_1}, \ \textit{T}), \\ [\text{conditions}] \end{array}$

 $\begin{array}{ll} \textbf{terminatedAt}(\textit{CE}, ~ \textit{T}) \leftarrow \\ \textbf{happensAt}(\textit{E}_{\textit{T}_{j}}, ~ \textit{T}), \\ [\text{conditions}] \end{array}$

where

conditions: ${}^{0-K}$ happensAt (E_k, T) , ${}^{0-M}$ holdsAt (F_m, T) , ${}^{0-N}$ atemporal-constraint

CE definition:

. . .

initiatedAt(CE, T) \leftarrow happensAt(E_{ln_1} , T), [conditions]

initiatedAt(CE, T) \leftarrow happensAt(E_{In_i} , T), [conditions] terminatedAt(CE, T) \leftarrow happensAt(E_{T_1} , T), [conditions]

terminatedAt(CE, T) \leftarrow happensAt(E_{T_j} , T), [conditions]

CE recognition:



CE definition:

. . .

initiatedAt(CE, T) \leftarrow happensAt(E_{ln_1} , T), [conditions]

initiatedAt(CE, T) \leftarrow happensAt(E_{In_i}, T), [conditions] terminatedAt(CE, T) \leftarrow happensAt(E_{T_1} , T), [conditions]

terminatedAt(CE, T) \leftarrow happensAt(E_{T_j} , T), [conditions]

CE recognition:



CE definition:

. . .

initiatedAt(CE, T) \leftarrow happensAt(E_{ln_1} , T), [conditions]

initiatedAt(CE, T) \leftarrow happensAt(E_{In_i} , T), [conditions] terminatedAt(CE, T) \leftarrow happensAt(E_{T_1} , T), [conditions]

terminatedAt(CE, T) \leftarrow happensAt(E_{T_j} , T), [conditions]

CE recognition:



CE definition:

initiatedAt(CE, T) \leftarrow happensAt(E_{In_1} , T), [conditions]

initiatedAt(CE, T) \leftarrow happensAt(E_{In_i} , T), [conditions] terminatedAt(CE, T) \leftarrow happensAt(E_{T_1} , T), [conditions]

 $\begin{array}{ll} \textbf{terminatedAt}(CE, \ T) \leftarrow \\ \textbf{happensAt}(E_{T_j}, \ T), \\ [\text{conditions}] \end{array}$

CE recognition: **holdsFor**(CE, I)



CE definition:

initiatedAt($leaving_object(P, Obj) = true, T$) \leftarrow happensAt(appear(Obj), T), holdsAt(inactive(Obj) = true, T), holdsAt(close(P, Obj) = true, T), holdsAt(person(P) = true, T) terminatedAt($leaving_object(P, Obj) = true, T$) \leftarrow happensAt(disappear(Obj), T)

CE recognition: holdsFor($leaving_object(P, Obj) = true, I$)

CE Definitions in the Run-Time Event Calculus: Statically Determined Fluents

> holdsFor(CE, I) \leftarrow holdsFor(F_1 , I_{F_1}), ..., holdsFor(F_f , I_{F_f}), interval_manipulation₁(I_{α} ,..., I_{ω}), ..., interval_manipulation_k(I_{A} ,..., I_{Ω})

where

```
interval\_manipulation(I_1, \dots, I_n, I):
union([I_1, \dots, I_n], I)
intersection([I_1, \dots, I_n], I)
relative\_complement(I_1, [I_2, \dots, I_n], I)
```

Interval Manipulation: Union



Interval Manipulation: Intersection



Interval Manipulation: Relative Complement



CE Definitions in the Run-Time Event Calculus: Statically Determined Fluents

CE definition:

 $\begin{aligned} & \mathsf{holdsFor}(abnormal(Vessel) = \mathsf{true}, \ l) \leftarrow \\ & \mathsf{holdsFor}(slowMotion(Vessel) = \mathsf{true}, \ l_1), \\ & \mathsf{holdsFor}(gap(Vessel) = \mathsf{true}, \ l_2), \\ & \mathsf{holdsFor}(stop(Vessel) = \mathsf{true}, \ l_3), \\ & \mathsf{union}([l_1, l_2, l_3], \ l) \end{aligned}$

CE Definitions in the Run-Time Event Calculus: Statically Determined Fluents

CE definition:

 $\begin{array}{ll} \mathsf{holdsFor}(abnormal(Vessel) = \mathsf{true}, \ l) \leftarrow \\ \mathsf{holdsFor}(slowMotion(Vessel) = \mathsf{true}, \ l_1), \\ \mathsf{holdsFor}(gap(Vessel) = \mathsf{true}, \ l_2), \\ \mathsf{holdsFor}(stop(Vessel) = \mathsf{true}, \ l_3), \\ \mathsf{union}([l_1, l_2, l_3], \ l) \end{array}$

Shorthand:

abnormal(Vessel) iff slowMotion(Vessel) or gap(Vessel) or idle(Vessel) CE Definitions in the Run-Time Event Calculus: Statically Determined Fluents Shorthand:

suspicious(Vessel1, Vessel2) iff
abnormal(Vessel1),
abnormal(Vessel2),
close(Vessel1, Vessel2),
not (in(Vessel1, Area) or in(Vessel2, Area))

CE Definitions in the Run-Time Event Calculus: Statically Determined Fluents Shorthand:

> suspicious(Vessel₁, Vessel₂) iff abnormal(Vessel₁), abnormal(Vessel₂), close(Vessel₁, Vessel₂), not (in(Vessel₁, Area) or in(Vessel₂, Area))

CE definition:

 $\begin{array}{l} \mathsf{holdsFor}(suspicious(Vessel_1, Vessel_2) = \mathsf{true}, \ I) \leftarrow \\ \mathsf{holdsFor}(abnormal(Vessel_1) = \mathsf{true}, \ I_1), \\ \mathsf{holdsFor}(abnormal(Vessel_2) = \mathsf{true}, \ I_2), \\ \mathsf{holdsFor}(close(Vessel_1, Vessel_2) = \mathsf{true}, \ I_2), \\ \mathsf{holdsFor}(in(Vessel_1, Area) = \mathsf{true}, \ I_4), \\ \mathsf{holdsFor}(in(Vessel_2, Area) = \mathsf{true}, \ I_5), \\ intersection([I_1, I_2, I_3], \ I_6), \\ relative_complement(I_6, \ [I_4, I_5], \ I) \\ \end{array}$

CE Hierarchies



CE Hierarchies











Run-Time Event Recognition

Real-time decision-support in the presence of:

- Very large SDE streams.
- Non-sorted SDE streams.
- SDE revision.
- Very large CE numbers.

Run-Time Event Calculus: Windowing



Run-Time Event Calculus: Windowing



Run-Time Event Calculus: Windowing



- Succinct representation \rightarrow code maintenance.
- ► Intuitive representation → domain experts unfamiliar with programming into the loop.

- \blacktriangleright Succinct representation \rightarrow code maintenance.
- \blacktriangleright Intuitive representation \rightarrow domain experts unfamiliar with programming into the loop.
- The full power of logic programming is available.
 - Complex atemporal computations.
 - Combination of events streams with static knowledge.

- \blacktriangleright Succinct representation \rightarrow code maintenance.
- \blacktriangleright Intuitive representation \rightarrow domain experts unfamiliar with programming into the loop.
- The full power of logic programming is available.
 - Complex atemporal computations.
 - Combination of events streams with static knowledge.
- Very efficient reasoning.
 - Even when event streams arrive with a delay.
 - Even in the presence of large specifications.

- \blacktriangleright Succinct representation \rightarrow code maintenance.
- ► Intuitive representation → domain experts unfamiliar with programming into the loop.
- The full power of logic programming is available.
 - Complex atemporal computations.
 - Combination of events streams with static knowledge.
- Very efficient reasoning.
 - Even when event streams arrive with a delay.
 - Even in the presence of large specifications.
- ► But:
 - The Event Calculus does not deal with uncertainty.

Event Recognition

Requirements:

- Efficient reasoning
 - to support real-time decision-making in large-scale, (geographically) distributed applications.
- Reasoning under uncertainty
 - to deal with various types of noise.
- Automated knowledge construction
 - to avoid the time-consuming, error-prone manual CE definition development.

Common Problems of Event Recognition

- Limited dictionary of SDE and context variables.
 - No explicit representation of oil spillage.
Common Problems of Event Recognition

- Limited dictionary of SDE and context variables.
 - ► No explicit representation of oil spillage.
- Incomplete SDE stream.
 - Sharp turn was not detected.

Common Problems of Event Recognition

- Limited dictionary of SDE and context variables.
 - No explicit representation of oil spillage.
- Incomplete SDE stream.
 - Sharp turn was not detected.
- Erroneous SDE detection.
 - Slow motion was classified as stop.

Common Problems of Event Recognition

- Limited dictionary of SDE and context variables.
 - No explicit representation of oil spillage.
- Incomplete SDE stream.
 - Sharp turn was not detected.
- Erroneous SDE detection.
 - Slow motion was classified as stop.
- Inconsistent ground truth (CE & SDE annotation).
 - Disagreement between (human) annotators.

Therefore, an adequate treatment of uncertainty is required.

Statistical Relational Learning



ProbLog

- ► A probabilistic logic programming language.
- Allows for independent probabilistic facts prob::fact.
 - prob indicates the probability that fact is part of a possible world.

ProbLog

- ► A probabilistic logic programming language.
- Allows for independent probabilistic facts prob::fact.
 - prob indicates the probability that fact is part of a possible world.
- Rules are written as in classic Prolog.

ProbLog

- ► A probabilistic logic programming language.
- Allows for independent probabilistic facts prob::fact.
 - prob indicates the probability that fact is part of a possible world.
- Rules are written as in classic Prolog.
- The probability of a query q imposed on a ProbLog database (success probability) is computed by the following formula:

$$P_s(q) = P(\bigvee_{e \in Proofs(q)} \bigwedge_{f_i \in e} f_i)$$

Event Recognition using ProbLog

Input	Output
340 0.45 :: <i>inactive</i> (<i>id</i> ₀)	340 0.41 :: $left_object(id_1, id_0)$
340 0.80 :: $p(id_0) = (20.88, -11.90)$	$340 \ 0.55 :: moving(id_2, id_3)$
340 0.55 :: <i>appear(id</i> ₀)	
340 0.15 :: <i>walking(id</i> ₂)	
340 0.80 :: $p(id_2) = (25.88, -19.80)$	
340 0.25 :: active(id ₁)	
340 0.66 :: $p(id_1) = (20.88, -11.90)$	
340 0.70 :: walking(id ₃)	
340 0.46 :: $p(id_3) = (24.78, -18.77)$	

Event Calculus in ProbLog



Event Calculus in ProbLog



Markov Logic Networks (MLN)

SYNTAX: weighted first-order logic formulas (w_i, F_i)

When input events SDE_A and SDE_B occur at T, then the output event *CE* is initiated:

3.18 happensAt(SDE_A, T) \land happensAt(SDE_B, T) \Rightarrow initiatedAt(CE, T)

Markov Logic Networks (MLN)

SYNTAX: weighted first-order logic formulas (w_i, F_i)

When input events SDE_A and SDE_B occur at T, then the output event *CE* is initiated:

3.18 happensAt(SDE_A, T) \land happensAt(SDE_B, T) \Rightarrow initiatedAt(CE, T)

SEMANTICS: (w_i, F_i) represents a probability distribution over possible worlds



A world violating formulas becomes less probable, but not impossible!

Event Calculus in Markov Logic Networks (MLN-EC)



MLN-EC: Probabilistic Inference

Marginal Inference:

 For all time points T, calculate the probability of each CE being true (recognised), given all input SDEs (evidence)

P(holdsAt(CE, T) = true|SDEs)

- Marginal inference is #P-complete \rightarrow approximate inference
- MC-SAT algorithm (Markov Chain Monte Carlo techniques with SAT solver)

MLN-EC: Probabilistic Inference

Maximum a Posteriori (MAP) Inference:

- Find the world with the highest probability
- Input: truth values for all input SDEs (evidence)
- Output: truth values of the output CEs that maximise the probability (recognition)

$$\operatorname{argmax}_{\operatorname{holdsAt}(CE, T)} \left(P(\operatorname{holdsAt}(CE, T)|SDEs) \right)$$

- MAP Inference is NP-hard \rightarrow approximate inference
- ► Various methods: local search, linear programming, etc.



∞ holdsAt(CE, T+1) \Leftarrow [Initiation Conditions]

 ∞ ¬holdsAt(CE, T+1) ¬holdsAt(CE), T) ∧ ¬[Initiation Conditions] ∞ ¬holdsAt(CE, T+1) ← [Termination Conditions]

∞ holdsAt(CE, T+1) \Leftarrow holdsAt(CE, T) \land \neg [Termination Conditions]



1.2 holdsAt(CE, T+1)⇐
[Initiation Conditions]

 ∞ ¬holdsAt(CE, T+1) ← ¬holdsAt(CE), T) ∧ ¬[Initiation Conditions]

∞ holdsAt(CE, T+1) \Leftarrow holdsAt(CE, T) \land \neg [Termination Conditions]



- 1.2 holdsAt(CE, T+1) \leftarrow [Initiation Conditions]
- 2.3 $\neg holdsAt(CE, T+1) \Leftarrow \\ \neg holdsAt(CE), T) \land \\ \neg [Initiation Conditions]$

- ∞ holdsAt(CE, T+1) \Leftarrow holdsAt(CE, T) \land \neg [Termination Conditions]



1.2 holdsAt(CE, T+1) \leftarrow [Initiation Conditions]

 ∞ ¬holdsAt(CE, T+1) ← ¬holdsAt(CE), T) ∧ ¬[Initiation Conditions]

2.3 holdsAt(CE, T+1) \Leftarrow holdsAt(CE, T) \land \neg [Termination Conditions]



1.2 holdsAt(CE, T+1)⇐
[Initiation Conditions]

 ∞ ¬holdsAt(CE, T+1) ← ¬holdsAt(CE), T) ∧ ¬[Initiation Conditions]

 $\begin{array}{ll} \textbf{0.6} & \text{holdsAt}(\text{CE}, \ \text{T}+1) \Leftarrow \\ & \text{holdsAt}(\text{CE}, \ \text{T}) \ \land \\ & \neg \begin{bmatrix} \text{Termination Conditions} \end{bmatrix} \end{array}$

MLN-EC: Complete Data



- Crisp Event Calculus
- Linear-chain Conditional Random Field
- Hard-constrained inertia
- Soft-constrained termination inertial rules
- Soft-constrained inertial rules

MLN-EC: Incomplete Data



	ECcrisp	:	Crisp Event Calculus
_	1-CRF	:	Linear-chain Conditional Random Field
-	$MLN-EC_{SIT}$	1	Soft-constrained termination inertial rules

MLN-EC: Summary

- We can deal with both:
 - Uncertainty in the CE definitions, and
 - uncertainty in the input data streams.
- Customisable inertia behaviour to meet the requirements of different applications.

MLN-EC: Summary

- We can deal with both:
 - Uncertainty in the CE definitions, and
 - uncertainty in the input data streams.
- Customisable inertia behaviour to meet the requirements of different applications.

But:

► There is room for improvement with respect to efficiency.

Event Recognition

Requirements:

- Efficient reasoning
 - to support real-time decision-making in large-scale, (geographically) distributed applications.
- Reasoning under uncertainty
 - to deal with various types of noise.
- Automated knowledge construction
 - ► to avoid the time-consuming, error-prone manual CE definition development.

Manual development of CE definitions:

- Time consuming.
- Error-prone.

Manual development of CE definitions:

- Time consuming.
- Error-prone.

- Learn complex definitions
 - Structure learning.

Manual development of CE definitions:

- Time consuming.
- Error-prone.

- Learn complex definitions
 - Structure learning.
- Learn from noisy data
 - Weight learning.

Manual development of CE definitions:

- Time consuming.
- Error-prone.

- Learn complex definitions
 - Structure learning.
- Learn from noisy data
 - Weight learning.
- Learn with incomplete or missing annotation
 - Semi-supervised, unsupervised learning.

Manual development of CE definitions:

- Time consuming.
- Error-prone.

- Learn complex definitions
 - Structure learning.
- Learn from noisy data
 - Weight learning.
- Learn with incomplete or missing annotation
 - Semi-supervised, unsupervised learning.
- Learn from large amounts of (streaming) data
 - Incremental learning, online learning.



Structure Learning with ILP

Inductive Logic Programming (ILP):

- Input: SDE streams annotated with CE
 - Positive (E^+) and negative (E^-) examples.
- Event recognition engine
 - Background knowledge B.
- Syntax of event recognition language
 - Language bias M.
- Output: A set of CE definitions
 - A hypothesis \mathcal{H} in the language of M, such that:

```
\underset{e^+ \in E^+}{\text{maximize}} B \cup \mathcal{H} \vDash e^+ \text{ and } \underset{e^- \in E^-}{\text{minimize}} B \cup \mathcal{H} \vDash e^-
```

A Generic ILP Algorithm



Online Learning



initiatedAt(moving(X, Y), T) \leftarrow happensAt(walking(X), T), happensAt(walking(Y), T), holdsAt(close(X, Y, 34), T).

Online Learning

Goal:

• Revise H to an H' that accounts for (all) the examples.



 $H': \begin{array}{c} \text{initiatedAt}(moving(X, Y), T) \leftarrow \\ happensAt(walking(X), T), \\ happensAt(walking(Y), T), \\ holdsAt(close(X, Y, 34), T), \\ holdsAt(orientation(X, Y, 45), T). \end{array}$

terminatedAt(moving(X, Y), T) \leftarrow happensAt(inactive(X), T), not holdsAt(close(X, Y, 34), T).

Online Learning


















We must start all over again

Markov Logic Networks: Online Structure Learning using background knowledge Axiomatization ($OSL\alpha$)



Training Example (Micro-Batch)

Simple Derived Events	Complex Event Annotation
$\texttt{HappensAt}(\texttt{walking}(\texttt{ID}_1), 99)$	
$\texttt{HappensAt}(\texttt{walking}(\texttt{ID}_2), 99)$	
$OrientationMove(ID_1, ID_2, 99)$	$HoldsAt(move(ID_1, ID_2), 100)$
$Close(ID_1, ID_2, 34, 99)$	
Next(99, 100)	

Hypergraph



Hypergraph and Relational Pathfinding

$$\mathbf{y}_{t}^{P} = \left(\neg \texttt{HoldsAt}(\texttt{MoveID}_{1}\texttt{ID}_{2}, 100) \right)$$

{HoldsAt(MoveID₁ID₂, 100),



Relational Pathfinding

$$\mathbf{y}_{t}^{P} = \left(\neg \texttt{HoldsAt}(\texttt{MoveID}_{1}\texttt{ID}_{2}, 100) \right)$$

{HoldsAt(MoveID₁ID₂, 100), Next(99, 100),



Relational Pathfinding

$$\mathbf{y}_t^P = \left(\neg \texttt{HoldsAt}(\texttt{MoveID}_1\texttt{ID}_2, 100)\right)$$

 $\label{eq:holdsAt(MoveID_1ID_2, 100), Next(99, 100),} $$ HappensAt(WalkingID_1, 99), $$$



Relational Pathfinding

$$\mathbf{y}_t^P = \left(\neg \texttt{HoldsAt}(\texttt{MoveID}_1\texttt{ID}_2, 100)\right)$$

 $\{ \label{eq:holdsAt} $$ {HoldsAt}(MoveID_1ID_2, 100), Next(99, 100), $$ HappensAt}(WalkingID_1, 99), $$ $$ HappensAt}(WalkingID_2, 99) $$ $$



Template-Guided Search

Wrongly predicted atom: ¬HoldsAt(MoveID₁ID₂, 100)



 $\begin{array}{l} \texttt{HoldsAt}(f,t_1) \Leftarrow \\ \texttt{InitiatedAt}(f,t_0) \land \\ \texttt{Next}(t_0,t_1) \end{array}$

Template-Guided Search

Wrongly predicted atom: ¬HoldsAt(MoveID₁ID₂, 100)



 $\begin{array}{l} \texttt{HoldsAt}(f,t_1) \Leftarrow \\ \texttt{InitiatedAt}(f,t_0) \land \\ \texttt{Next}(t_0,t_1) \end{array}$

 $\begin{array}{l} \texttt{HoldsAt}(\texttt{MoveID}_1\texttt{ID}_2, \texttt{100}) \Leftarrow \\ \texttt{InitiatedAt}(\texttt{MoveID}_1\texttt{ID}_2, t_0) \land \\ \texttt{Next}(t_0, \texttt{100}) \end{array}$

Template-Guided Search

Wrongly predicted atom: ¬HoldsAt(MoveID₁ID₂, 100)



 $\begin{array}{l} \texttt{HoldsAt}(f,t_1) \Leftarrow \\ \texttt{InitiatedAt}(f,t_0) \land \\ \texttt{Next}(t_0,t_1) \end{array}$

 $\begin{array}{l} \texttt{HoldsAt}(\texttt{MoveID}_1\texttt{ID}_2, \texttt{100}) \Leftarrow \\ \texttt{InitiatedAt}(\texttt{MoveID}_1\texttt{ID}_2, t_0) \land \\ \texttt{Next}(t_0, \texttt{100}) \end{array}$



InitiatedAt(MoveID₁ID₂, 99)

Reduced Hypergraph



Clause Creation, Clause Evaluation and Weight Learning

Clause creation:

Generalize each path into a definite clause

 $\texttt{InitiatedAt}(\texttt{move}(\textit{id}_1, \textit{id}_2), t) \Leftarrow \\ \texttt{HappensAt}(\texttt{walking}(\textit{id}_1), t) \land \\ \texttt{HappensAt}(\texttt{walking}(\textit{id}_2), t)$

Clause Creation, Clause Evaluation and Weight Learning

Clause creation:

Generalize each path into a definite clause

 $\texttt{InitiatedAt}(\texttt{move}(\textit{id}_1, \textit{id}_2), t) \Leftarrow \\ \texttt{HappensAt}(\texttt{walking}(\textit{id}_1), t) \land \\ \texttt{HappensAt}(\texttt{walking}(\textit{id}_2), t)$

Clause evaluation:

Keep clauses whose coverage of the annotation is significantly greater than that of the clauses already learnt. Clause Creation, Clause Evaluation and Weight Learning

Clause creation:

Generalize each path into a definite clause

 $\texttt{InitiatedAt}(\texttt{move}(\textit{id}_1, \textit{id}_2), t) \Leftarrow \\ \texttt{HappensAt}(\texttt{walking}(\textit{id}_1), t) \land \\ \texttt{HappensAt}(\texttt{walking}(\textit{id}_2), t)$

Clause evaluation:

Keep clauses whose coverage of the annotation is significantly greater than that of the clauses already learnt.

Weight learning:

- Extended clauses inherit initially the weights of their ancestors.
- Optimize the weights of all clauses.

Machine Learning for Event Recognition: Summary

- CE definition construction
 - using large (streaming) data;
 - tolerant to noise.
- Acceptable predictive accuracy.

Machine Learning for Event Recognition: Summary

- CE definition construction
 - using large (streaming) data;
 - tolerant to noise.
- Acceptable predictive accuracy.
- But:
 - Constructed CE definitions tend to overfit the data.
 - There is room for improvement with respect to efficiency.

Open Issues

Issue (1): Real-time Probabilistic Event Recognition

Types of uncertainty:

- Noisy input stream.
- Imprecise CE definitions.

Existing solution:

Probabilistic graphical models & logic-based approaches.

Issue (1): Real-time Probabilistic Event Recognition

Types of uncertainty:

- Noisy input stream.
- Imprecise CE definitions.

Existing solution:

Probabilistic graphical models & logic-based approaches.

Major drawback:

Large overhead that does not allow for real-time performance.

Issue (2): Distributed Recognition

Distribution of CE recognition among multiple nodes:

- Data distribution.
- Distribution of CE (sub-)definitions.

Issue (2): Distributed Recognition

Distribution of CE recognition among multiple nodes:

- Data distribution.
- Distribution of CE (sub-)definitions.

But, virtually no work on distributing probabilistic CE recognition.

Issue (2): Distributed Recognition

Distribution of CE recognition among multiple nodes:

- Data distribution.
- Distribution of CE (sub-)definitions.

But, virtually no work on distributing probabilistic CE recognition.

Interplay with communication efficiency:

- Reduce communication by decomposing recognition in set of local constraints at event sources
 - By sketching.
 - By geometric models.
- Limited to particular types of CE: functions over aggregate values.

Issue (3): Multi-scale Temporal Aggregation of Events

CE evolve over multiple scales of time and space:

- SDE streams.
- Context information, e.g. historic data.



Issue (3): Multi-scale Temporal Aggregation of Events

CE evolve over multiple scales of time and space:

- SDE streams.
- Context information, e.g. historic data.

The recognition system should be adaptable, computing dynamically the appropriate scales

- Appropriate lengths of multi-granular windows.
- Appropriate slice of context information.



Issue (3): Multi-scale Temporal Aggregation of Events

CE evolve over multiple scales of time and space:

- SDE streams.
- Context information, e.g. historic data.

The recognition system should be adaptable, computing dynamically the appropriate scales

- Appropriate lengths of multi-granular windows.
- Appropriate slice of context information.

Semantics and processing guarantees, despite adaptation.



Issue (4): Event Forecasting

Current work focuses on

- Present (monitoring).
- Past (post mortem analysis).



Issue (4): Event Forecasting

Current work focuses on

- Present (monitoring).
- Past (post mortem analysis).

Stream processing will focus on future

- Predictive analysis.
- Trend detection.



Issue (4): Event Forecasting

Current work focuses on

- Present (monitoring).
- Past (post mortem analysis).

Stream processing will focus on future

- Predictive analysis.
- Trend detection.

The earlier the better

- Time to react to an anticipated situation.
- Avoidance of situation or initialisation of counter measures.





Resources

- Tutorial paper: A. Artikis, A. Skarlatidis, F. Portet, G. Paliouras: Logic-based event recognition. Knowledge Engineering Review 27(4): 469-506 (2012)
- Slides, complex event recognition software, datasets: http://cer.iit.demokritos.gr


Acknowledgements

- Marek Sergot, Imperial College London
- Anastasios Skarlatidis, Pollfish
- Matthias Weidlich, Humboldt-Universität zu Berlin