# Complex Event Recognition with Allen Relations 

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http://cer.iit.demokritos.gr/


# Complex Event Recognition 



## Complex Event Recognition


https://cer.iit.demokritos.gr (maritime)

## Event Calculus

- A logic programming language for representing and reasoning about events and their effects.
- Key components:
- event (typically instantaneous).
- fluent: a property that may have different values at different points in time.

Robert A. Kowalski, Marek J. Sergot: A Logic-based Calculus of Events. New Gener. Comput. 4(1): 67-95, 1986.

## Event Calculus

- A logic programming language for representing and reasoning about events and their effects.
- Key components:
- event (typically instantaneous).
- fluent: a property that may have different values at different points in time.
- Built-in representation of inertia:
- $F=V$ holds at a particular time-point if $F=V$ has been initiated by an event at some earlier time-point, and not terminated by another event in the meantime.

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## Run-Time Event Calculus (RTEC): <br> Fluent Specification

Simple Fluents:
initiatedAt $(F=V, T) \leftarrow$
happensAt $\left(E_{I_{1}}, T\right)[$,
conditions].
terminated $\mathbf{A t}(F=V, T) \leftarrow$
happensAt $\left(E_{T_{1}}, T\right)[$,
conditions].
where conditions:
${ }^{0-K}$ [not] happensAt $\left(E_{k}, T\right)$,
$0-M$ [not] holdsAt $\left(F_{m}=V_{m}, T\right)$,
$0-N$ atemporal-constraint ${ }_{n}$
Artikis A., Sergot M. and Paliouras G., An Event Calculus for Event Recognition. In IEEE Transactions on Knowledge and Data Engineering (TKDE), 27(4), 895-908, 2015.

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Statically Determined Fluents:
holdsFor $(F=V, I) \leftarrow$
holdsFor $\left(F_{1}=V_{1}, I_{1}\right)[$,
holdsFor $\left(F_{2}=V_{2}, I_{2}\right), \ldots$
holdsFor $\left(F_{n}=V_{n}, I_{n}\right)$,
intervalConstruct $\left(L_{1}, I_{n+1}\right), \ldots$
intervalConstruct $\left.\left(L_{m}, I\right)\right]$.
where intervalConstruct:
union_all or
intersect_all or
relative_complement_all

## Statically Determined Fluent: Anchored or Moored

holdsFor(anchoredOrMoored (Vessel) $=$ true, $I) \leftarrow$ holdsFor (stopped (Vessel) $=$ farFromPorts, $I_{s f}$ ), holdsFor (withinArea(Vessel, anchorage) $=$ true, $I_{\text {wa }}$ ), intersect_all $\left(\left[I_{s f}, I_{w a}\right], I_{s a}\right)$, holdsFor(stopped $($ Vessel $)=$ nearPorts, $\left.I_{s n}\right)$, union_all( $\left.\left[I_{s a}, I_{s n}\right], I\right)$.

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## Interval Constructs \& Allen Relations



## RTEC $_{A}$ : RTEC with Allen Relations

holdsFor(disappearedInArea(Vessel, Area Type) $=$ true, $I) \leftarrow$
holdsFor (withinArea(Vessel, AreaType) = true, $\mathcal{S}$ ), holdsFor (gap (Vessel) $=$ farFromPorts, $\mathcal{T})$, allen(meets, $\mathcal{S}, \mathcal{T}$, target, $I)$.

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Query time: $q_{82}$


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## RTEC $_{A}$ : Correctness \& Complexity

## Correctness of RTEC $A$

RTEC ${ }_{A}$ computes all maximal intervals of a fluent defined in terms of an Allen relation, and no other interval.

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## Complexity of RTEC $A$

The cost of computing the maximal intervals of a fluent defined in terms of an Allen relation is $\mathcal{O}(n)$, where $n$ is the number of input intervals.

## Experimental Evaluation



Code, Data \& Temporal Specifications:

- https://github.com/aartikis/RTEC/tree/allen


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## Summary \& Further Work

## RTEC $_{A}$ :

- An open-source complex event recognition framework.
- Support for Allen relations in event patterns.
- Correct Allen relation computation with windowing.
- Linear time complexity.
- Reproducible empirical evaluation on large, real data streams.


## Summary \& Further Work

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Further Work:

- Support approximate Allen relations.
- Contrast Allen relation with event sequencing operators.
- Support events with delayed effects.

Appendix

## Run-Time Event Calculus (RTEC)

## Predicate

happensAt $(E, T)$
initiatedAt $(F=V, T)$
terminatedAt $(F=V, T)$
holdsFor $(F=V, I)$
holdsAt $(F=V, T)$
union_all( $\left.\left[J_{1}, \ldots, J_{n}\right], I\right)$
intersect_all([ $\left.\left.J_{1}, \ldots, J_{n}\right], I\right)$
relative_complement_all $\left(I^{\prime},\left[J_{1}, \ldots, J_{n}\right], I\right)$

## Meaning

Event $E$ occurs at time $T$
At time $T$ a period of time for which $F=V$ is initiated

At time $T$ a period of time for which $F=V$ is terminated
$I$ is the list of the maximal intervals for which $F=V$ holds continuously

The value of fluent $F$ is $V$ at time $T$
$I=\left(J_{1} \cup \ldots \cup J_{n}\right)$
$I=\left(J_{1} \cap \ldots \cap J_{n}\right)$
$I=I^{\prime} \backslash\left(J_{1} \cup \ldots \cup J_{n}\right)$

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## Interval Manipulation: Relative Complement

relative_complement_all


## RTEC $_{A}$ : RTEC with Allen Relations

holdsFor(suspiciousRendezVous $\left(\right.$ Vessel $_{1}$, Vessel $\left.\left._{2}\right)=\operatorname{true}, I\right) \leftarrow$ holdsFor $\left(\operatorname{gap}\left(\operatorname{Vesse}_{1}\right)=\right.$ farFromPorts, $\left.I_{g_{1}}\right)$, holdsFor $\left(\operatorname{gap}\left(\right.\right.$ Vessel $\left._{2}\right)=$ farFromPorts, $\left.\mathrm{I}_{\mathrm{g}_{2}}\right)$,
holdsFor $\left(\right.$ proximity $\left(\right.$ Vessel $_{1}$, Vessel $\left.\left._{2}\right)=\operatorname{true}, \mathcal{T}\right)$, union_all( $\left.\left[I_{g_{1}}, l_{g_{2}}\right], \mathcal{S}\right)$, allen(during, $\mathcal{S}, \mathcal{T}$, target, $/$ ).


## Experimental Evaluation

> Batch setting.

| Win | ow size | Reasoni | Time | Outp Interva | ut Pairs |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Days | Input Intervals | RTEC ${ }_{\text {A }}$ | $\mathrm{D}^{2} \mathrm{I}$ A | $\mathrm{RTEC}_{A}$ | $\mathrm{D}^{2} \mathrm{I} A$ |
| 1 | 125 | 1 | 48 | 5K | 5 K |
| 2 | 250 | 2 | 164 | 19K | 18K |
| 4 | 500 | 4 | 568 | 72K | 71K |
| 8 | 1K | 8 | 1.7K | 237K | 236K |
| 16 | 2K | 15 | 7.8K | 878K | 874K |

