Scalable Complex Event Recognition

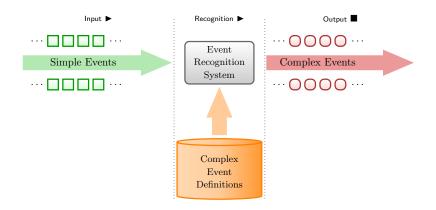
Efthimis Tsilionis

National and Kapodistrian University of Athens, Greece NCSR Demokritos, Greece





Complex Event Recognition



Logic-based approaches:

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Event Calculus: Reasoning about events and their effects.

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Temporal frameworks:

- ▶ Event Calculus: Reasoning about events and their effects.
- Event Calculus for Run-Time Reasoning (RTEC): EC with optimization techniques for CER.

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Motivation: Modify the *what* and *how* to achieve scalability. Contributions:

- ▶ $RTEC_{inc}$: Incremental RTEC (the *what*).
- tensor-EC: Tensor-based formalization of EC (the *how*).

Publications

Journal Publications:

- <u>Tsilionis E.</u>, Artikis A., Paliouras G., Incremental Event Calculus for Run-Time Reasoning. In *Journal of Artificial Intelligence Research (JAIR)*, 73, pp. 967—1023, 2022.
- <u>Tsilionis E.</u>, Koutroumanis N., Nikitopoulos P., Doulkeridis C. and Artikis A., Online Event Recognition from Moving Vehicles. In *Theory and Practice of Logic Programming (TPLP)*, 19(5-6), pp. 841—856, 2019.

Conference Publications:

- ► <u>Tsilionis E.</u>, Artikis A., Paliouras G., A Tensor-Based Formalization of the Event Calculus. In *Proceedings of the Thirty-Third International Joint Conference on Artificial Intelligence, IJCAI-24*, pp. 3584–3592, 2024.
- <u>Tsilionis E.</u>, Artikis A., Paliouras G., Incremental Event Calculus for Run-Time Reasoning (Extended Abstract). In *Proceedings of the Thirty-Second International Joint Conference on Artificial Intelligence, IJCAI-23*, pp. 6974-6978, 2023.
- ► <u>Tsilionis E.</u>, Artikis A. and Paliouras G., Incremental Event Calculus for Run-Time Reasoning. In 13th International Conference on Distributed and Event-Based Systems (DEBS), pp. 79–90, 2019.

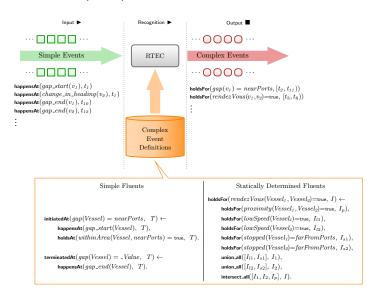
Background

Event Calculus

- ► A logic programming language for representing and reasoning about events and their effects.
- Key components:
 - event (typically instantaneous).
 - fluent: a property that may have different values at different points in time.
- ▶ Built-in representation of inertia:
 - ightharpoonup F=V holds at a particular time-point if F=V has been initiated by an event at some earlier time-point, and not terminated by another event in the meantime.

Background

Run-Time Event Calculus (RTEC)



```
\begin{split} & \mathbf{initiatedAt}(gap(Vessel) {=} farFromPorts, \ T) \leftarrow \\ & \mathbf{happensAt}(gap\_start(Vessel), \ T), \\ & \mathbf{not} \ \ \mathbf{holdsAt}(withinArea(Vessel, nearPorts) {=} \mathbf{true}, \ T). \end{split}
```

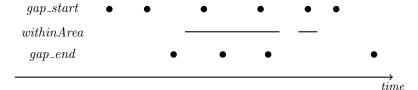
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Reasoning: holdsFor(gap(Vessel) = farFromPorts, I)

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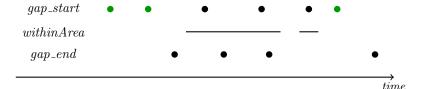
```
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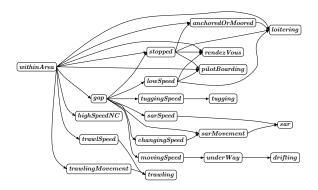


```
initiatedAt(qap(Vessel) = farFromPorts, T) \leftarrow
        happensAt(qap\_start(Vessel), T),
        not holdsAt(withinArea(Vessel, nearPorts)=true, T).
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        happensAt(qap\_end(Vessel), T).
  Reasoning: holdsFor(gap(Vessel) = farFromPorts, I)
qap\_start
withinArea
 qap\_end
```

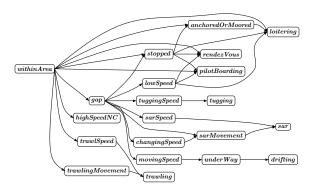
time

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initiatedAt(qap(Vessel) = farFromPorts, T) \leftarrow
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farFromPorts
  qap\_start
 withinArea
  qap\_end
                                                               time
```

Hierarchical Event Descriptions



Hierarchical Event Descriptions

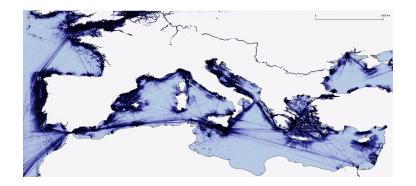


Semantics

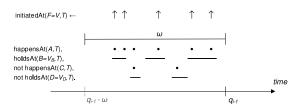
An event description is a locally stratified logic program, i.e., it has a **unique perfect** model.

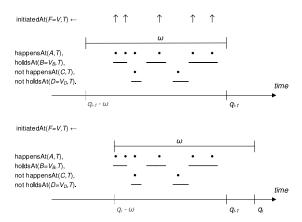
Part 1: Modify the what

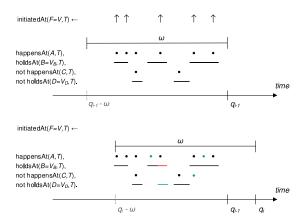
Incremental CER

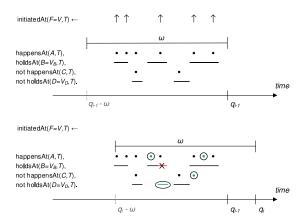


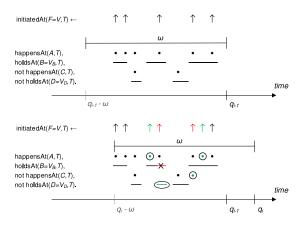
- ▶ Delayed events (e.g., satellite GPS messages)
- Overlapping temporal windows



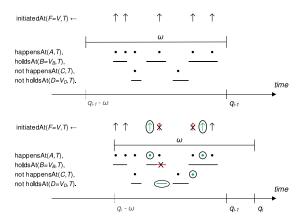




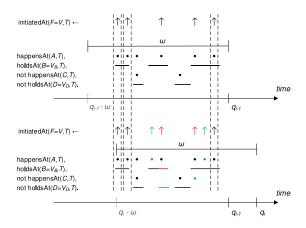


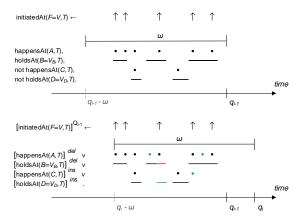


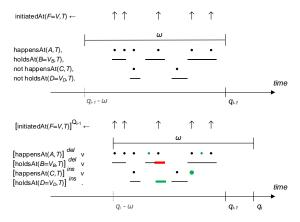
Problem Statement RTEC

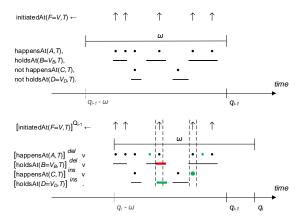


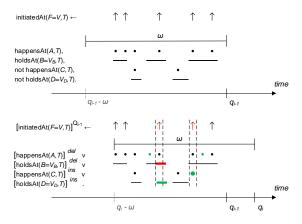
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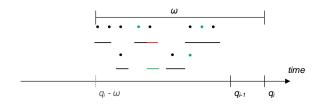


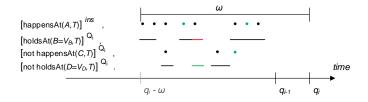


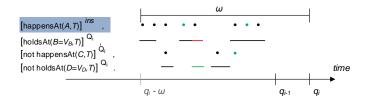


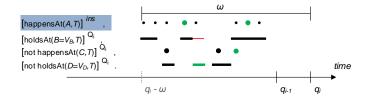


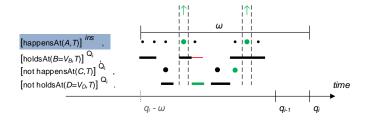
$$\begin{split} & \text{initiatedAt}(F=V,\,T) \leftarrow & \text{initiatedAt}(F=V,\,T) \leftarrow \\ & \left[\text{happensAt}(A,\,T) \right]^{ins}, & \left[\text{happensAt}(A,\,T) \right]^{Q_{\text{l}} \setminus ins}, \\ & \left[\text{holdsAt}(B=V_B,\,T) \right]^{Q_{\text{l}}}, & \text{(a)} & \left[\text{holdsAt}(B=V_B,\,T) \right]^{ins}, & \text{(b)} \\ & \text{not} \left[\text{happensAt}(C,\,T) \right]^{Q_{\text{l}}}, & \text{not} \left[\text{happensAt}(C,\,T) \right]^{Q_{\text{l}}}, \\ & \text{not} \left[\text{holdsAt}(D=V_D,\,T) \right]^{Q_{\text{l}}}, & \text{not} \left[\text{holdsAt}(D=V_D,\,T) \right]^{Q_{\text{l}}}. \\ & \text{initiatedAt}(F=V,\,T) \leftarrow & \text{initiatedAt}(F=V,\,T) \leftarrow \\ & \left[\text{happensAt}(A,\,T) \right]^{Q_{\text{l}} \setminus ins}, & \left[\text{happensAt}(A,\,T) \right]^{Q_{\text{l}} \setminus ins}, \\ & \left[\text{holdsAt}(B=V_B,\,T) \right]^{Q_{\text{l}} \setminus ins}, & \left[\text{holdsAt}(B=V_B,\,T) \right]^{Q_{\text{l}} \setminus ins}, \\ & \text{not} \left[\text{holdsAt}(B=V_B,\,T) \right]^{Q_{\text{l}} \setminus ins}, & \text{not} \left[\text{happensAt}(C,\,T) \right]^{Q_{\text{l}} \cup del}. \\ \end{split}$$











$RTEC_{inc}$: Formal properties

Correctness

 $RTEC_{inc}$ computes exactly the same intervals of the fluents of an event description as RTEC, and no other interval.

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Complexity

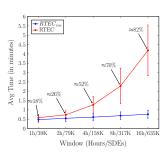
The most important factor for performance improvement is the ratio of delayed insertions/retractions to the degree of overlap:

$$\frac{n \times e}{m_{ov}} < 1$$

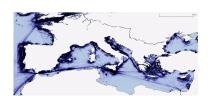
RTEC_{inc}: Empirical Evaluation

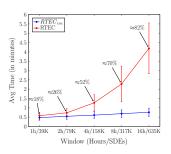
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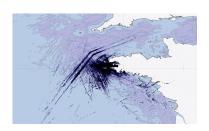


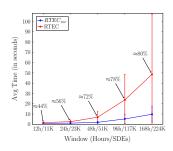


$RTEC_{inc}$: Empirical Evaluation









Logical Inference in Tensor Spaces

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- Logical reasoning through algebraic operations is a step towards neuro-symbolic integration.
- Use of efficient (parallel) algorithms and great computing resources (GPUs).

ightharpoonup Domain entities: N

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 $\mathsf{happensAt}(e(X,Y),\ T)$

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 $\mathsf{happensAt}(e(\textcolor{red}{X},\textcolor{red}{Y}),\ T)$

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 $\mathsf{happensAt}(e(\pmb{X},\pmb{Y}),\ T)$

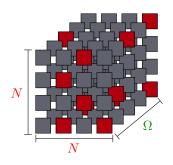
- ▶ Domain entities: *N*
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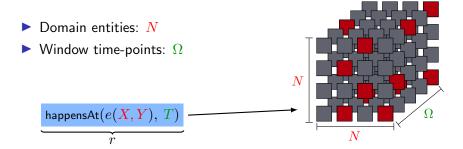
$$\underbrace{\frac{\mathsf{happensAt}(e(\pmb{X},\pmb{Y}),\,T)}{r}}$$

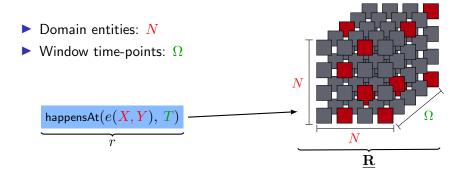
ightharpoonup Domain entities: N

ightharpoonup Window time-points: Ω

 $\underbrace{\frac{\mathsf{happensAt}(e(X\!\!\!X,Y),\ T)}{r}}$







 $\begin{array}{c} {\color{red} \triangleright} \ \, \text{Domain entities: } \, N \\ {\color{red} \triangleright} \ \, \text{Window time-points: } \, \Omega \\ \\ {\color{red} \underline{\hspace{0.5cm}}} \\ {\color{red} happensAt}(e(X,Y),T) \\ \\ {\color{red} r} \end{array}$

$$\underline{\mathbf{R}}_{i,j,k} = \begin{cases} 1, & \text{if } \mathbf{M}_P \models r, \text{ for } c_i, c_j, t_k \\ \\ 0, & \text{o.w} \end{cases}$$

$$\forall \ 1 \leq i, j \leq N, 1 \leq k \leq \Omega \ .$$

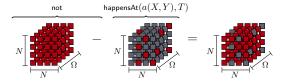
Tensor-EC: Reasoning

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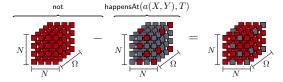
Negation:

Tensor-EC: Reasoning

Negation:

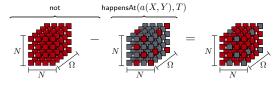


Negation:

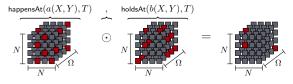


${\bf Conjunction:}$

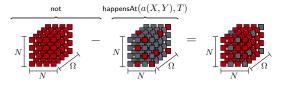
Negation:



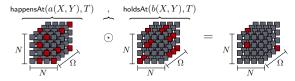
Conjunction:



Negation:

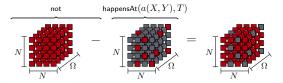


Conjunction:

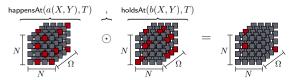


${\bf Disjunction:}$

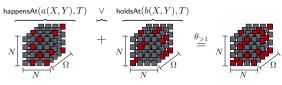
Negation:



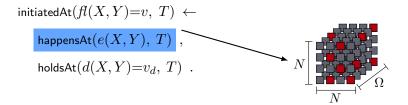
Conjunction:

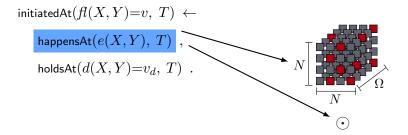


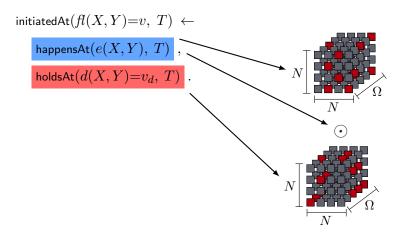
Disjunction:

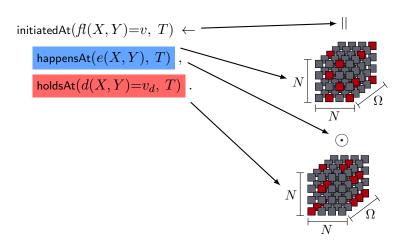


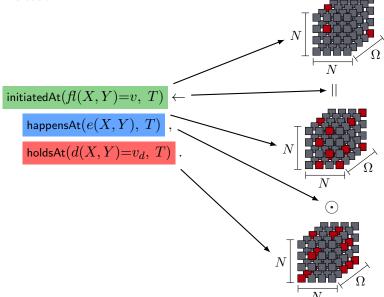
```
\begin{split} & \mathsf{initiatedAt}(f\!l(X,Y) {=} v,\ T) \ \leftarrow \\ & \mathsf{happensAt}(e(X,Y),\ T)\ , \\ & \mathsf{holdsAt}(d(X,Y) {=} v_d,\ T)\ . \end{split}
```











```
\begin{split} & \operatorname{holdsAt}(f\!l(X,Y)\!=\!v,\;T) \leftarrow & \operatorname{holdsAt}(f\!l(X,Y)\!=\!v,\;T) \leftarrow \\ & \operatorname{initiatedAt}(f\!l(X,Y)\!=\!v,\;T_{prev}), & \operatorname{holdsAt}(f\!l(X,Y)\!=\!v,\;T_{prev}), \\ & \operatorname{not\; terminatedAt}(f\!l(X,Y)\!=\!v,\;T_{prev}), & \operatorname{not\; terminatedAt}(f\!l(X,Y)\!=\!v,\;T_{prev}), \\ & \operatorname{next}(T_{prev},T). & \operatorname{next}(T_{prev},T). \end{split}
```

$\mathsf{holdsAt}(fl(X,Y){=}v,\ T) \leftarrow$

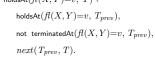
```
\begin{split} & \text{initiatedAt}(f\!l(X,Y){=}v,\ T_{prev}), \\ & \text{not terminatedAt}(f\!l(X,Y){=}v,\ T_{prev}), \\ & next(T_{prev},T). \end{split}
```

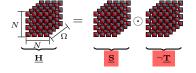
$\mathsf{holdsAt}(fl(X,Y){=}v,\ T) \leftarrow$

```
\begin{aligned} & \operatorname{holdsAt}(f\!l(X,Y) {=} v, \ T_{prev}), \\ & \operatorname{not \ terminatedAt}(f\!l(X,Y) {=} v, \ T_{prev}), \\ & next(T_{prev}, T). \end{aligned}
```

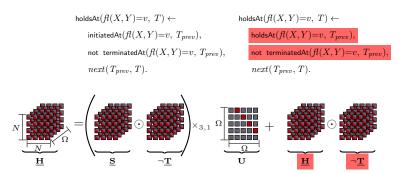


$$\begin{split} & \mathsf{holdsAt}(f\!l(X,Y)\!=\!v,\ T) \leftarrow & & \mathsf{holdsAt}(f\!l(X,Y)\!=\!v,\ T) \leftarrow \\ & & \mathsf{initiatedAt}(f\!l(X,Y)\!=\!v,\ T_{prev}), & & \mathsf{holdsAt}(f\!l(X,Y)\!=\!v,\ T_{pr}), \\ & & \mathsf{not\ terminatedAt}(f\!l(X,Y)\!=\!v,\ T_{prev}), & & \mathsf{not\ terminatedAt}(f\!l(X,Y)\!=\!v,\ T_{prev}), \\ & & & & \mathsf{next}(T_{prev},\ T). & & & & \mathsf{next}(T_{prev},\ T). \end{split}$$





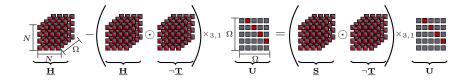
$$\begin{aligned} & \operatorname{holdsAt}(f\!l(X,Y)\!=\!v,\ T) \leftarrow & \operatorname{holdsAt}(f\!l(X,Y)\!=\!v,\ T) \leftarrow \\ & \operatorname{initiatedAt}(f\!l(X,Y)\!=\!v,\ T_{prev}), & \operatorname{holdsAt}(f\!l(X,Y)\!=\!v,\ T_{prev}), \\ & \operatorname{not\ terminatedAt}(f\!l(X,Y)\!=\!v,\ T_{prev}), & \operatorname{not\ terminatedAt}(f\!l(X,Y)\!=\!v,\ T_{prev}), \\ & \underbrace{\operatorname{next}(T_{prev},T).} & \operatorname{next}(T_{prev},T). \end{aligned}$$

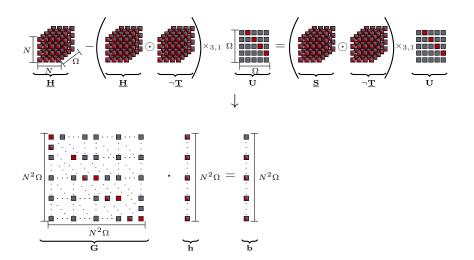


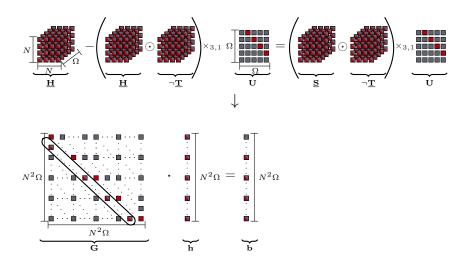
$$\begin{aligned} & \operatorname{holdsAt}(fl(X,Y) = v, \ T) \leftarrow & \operatorname{holdsAt}(fl(X,Y) = v, \ T) \leftarrow \\ & \operatorname{initiatedAt}(fl(X,Y) = v, \ T_{prev}), & \operatorname{holdsAt}(fl(X,Y) = v, \ T_{prev}), \\ & \operatorname{not \ terminatedAt}(fl(X,Y) = v, \ T_{prev}), & \operatorname{not \ terminatedAt}(fl(X,Y) = v, \ T_{prev}), \\ & \operatorname{next}(T_{prev}, T). & & \operatorname{next}(T_{prev}, T). \end{aligned}$$

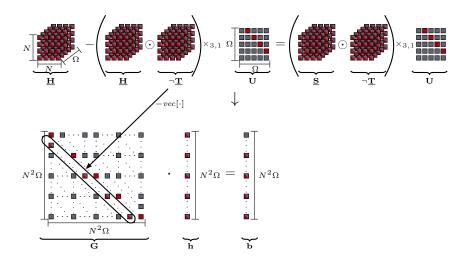
$$\text{holdsAt}(fl(X,Y) = v, \ T) \leftarrow \\ \text{initiatedAt}(fl(X,Y) = v, \ T_{prev}), \\ \text{not terminatedAt}(fl(X,Y) = v, \ T_{prev}), \\ \text{next}(T_{prev}, T). \\$$

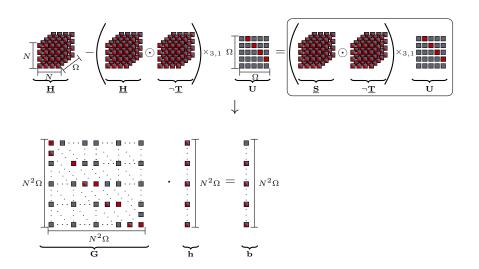
$$\begin{array}{c} \operatorname{holdsAt}(fl(X,Y) = v, \ T) \leftarrow \\ \operatorname{initiatedAt}(fl(X,Y) = v, \ T_{prev}), \\ \operatorname{not \ terminatedAt}(fl(X,Y) = v, \ T_{prev}), \\ \operatorname{not \ terminatedAt}(fl(X,Y) = v, \ T_{prev}), \\ \operatorname{next}(T_{prev}, T). \end{array} \\ \begin{array}{c} \operatorname{not \ terminatedAt}(fl(X,Y) = v, \ T_{prev}), \\ \operatorname{next}(T_{prev}, T). \end{array} \\ \\ N \\ \begin{array}{c} \underline{\underline{\mathbf{H}}} \\ \underline{\underline{\mathbf{N}}} \\ \underline{$$

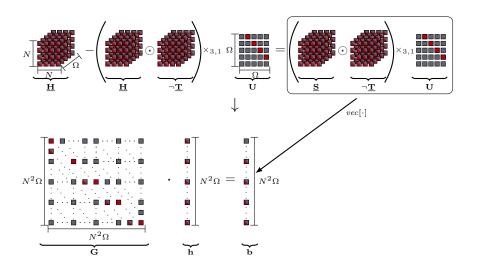


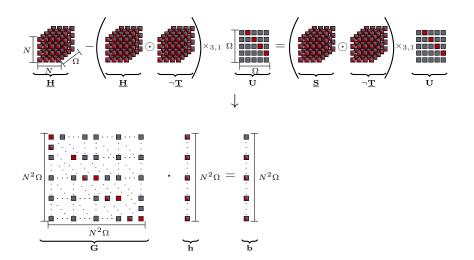


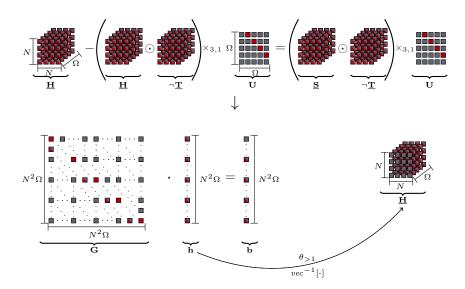












Tensor-EC: Formal Properties

Correctness

The **unique solution** of the equation coincides with the time-points at which a fluent-value pair holds, as expressed by the **perfect model** of the corresponding logic program.

Tensor-EC: Formal Properties

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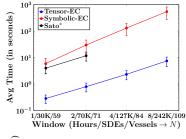
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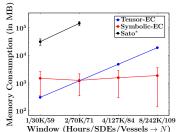
Complexity

The time complexity of solving the equation is $\mathcal{O}(N^{p-1}\Omega)$ for order-p tensors.

Tensor-EC: Experimental Evaluation

Brest

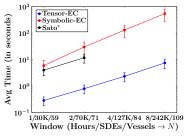


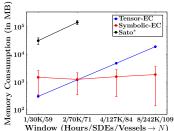


^{*}Sato, T. A linear algebraic approach to datalog evaluation. Theory and Practice of Logic Programming, 17(3):244–265, 2017. 22/24

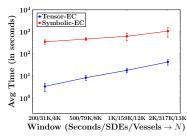
Tensor-EC: Experimental Evaluation

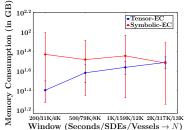
Brest





European seas





^{*}Sato, T. A linear algebraic approach to datalog evaluation. Theory and Practice of Logic Programming, 17(3):244-265, 2017. 22/24

Scalable Complex Event Recognition

▶ Modify the *what* is to be computed — $RTEC_{inc}^{\dagger}$:

[†]https://github.com/eftsilio/Incremental_RTEC

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 - ► Handles delays/revisions in the input.

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 - Performance improvement wrt state-of-the-art.

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- Probabilistic CER in tensor spaces.

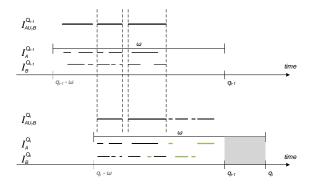
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- Probabilistic CER in tensor spaces.
- ► Neuro-symbolic CER.

Appendix

$RTEC_{inc}$: Statically Determined Fluents

Union

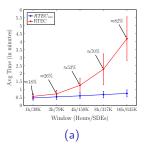
$$I_F^{\mathsf{Q}_i} = \left[(I_{A \cup_T B}^{\mathsf{Q}_{i-1}}) \setminus_T \left[\left[(I_A^{del} \cup_T I_B^{del}) \setminus_T (I_{A \cap_T B}^{\mathsf{Q}_{i-1}}) \right] \cup_T (I_A^{del} \cap_T I_B^{del}) \right] \right] \cup_T (I_A^{ins} \cup_T I_B^{ins})$$

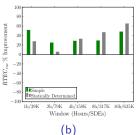


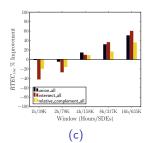
$RTEC_{inc}$: Evaluation (Natural Delays)

European seas

- Delays up to 16 hours
- ▶ 17M position signals, 34K vessels
- ▶ January 2016

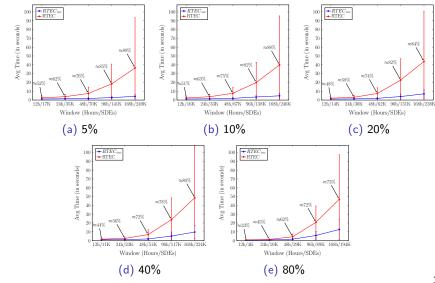






Brest

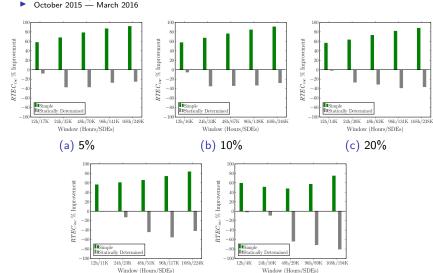
- 5M position signals, 5K vessels
- October 2015 March 2016



(d) 40%

Brest

5M position signals, 5K vessels

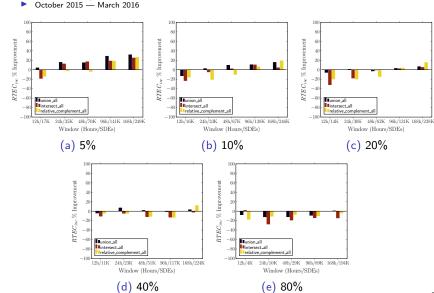


(e) 80%

4/7

Brest

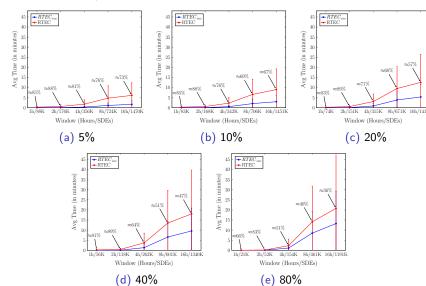
ightharpoonup 5M position signals, 5K vessels



5/7

Fleet Management

- ▶ 70M position signals, 6K vessels
- ▶ June 2018 August 2018



Tensor-EC: Evaluation

Simple Event Description

- One fluent and three input events.
- ► The fluent is defined by one initiatedAt and one terminatedAt rule, plus the *inertia* axiom.
- ► Temporal window of 20 time-points.

Method	Reasoning time (msec)	Memory (MB)
tensor-EC	1	0.01
symbolic-EC	0	5
Sato§	1	0.01
Sakama et al.¶	9205	475.3

Sato, T. A linear algebraic approach to datalog evaluation. Theory and Practice of Logic Programming, 17(3):244–265, 2017.

Sakama et al. Logic programming in tensor spaces. Annals of Mathematics and Artificial Intelligence, 89, 12 2021.